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# Security and Performance Implications of BGP Rerouting-resistant Guard Selection Algorithms for Tor



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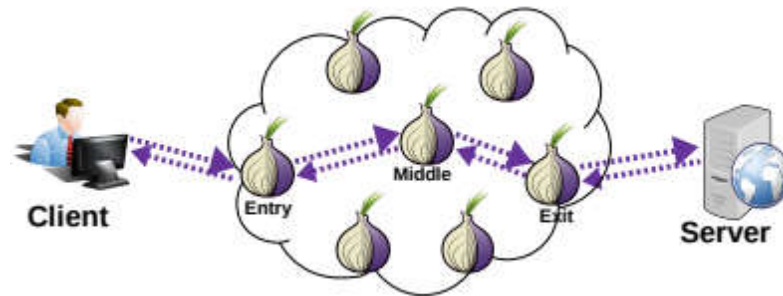


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# Motivation (1/2)

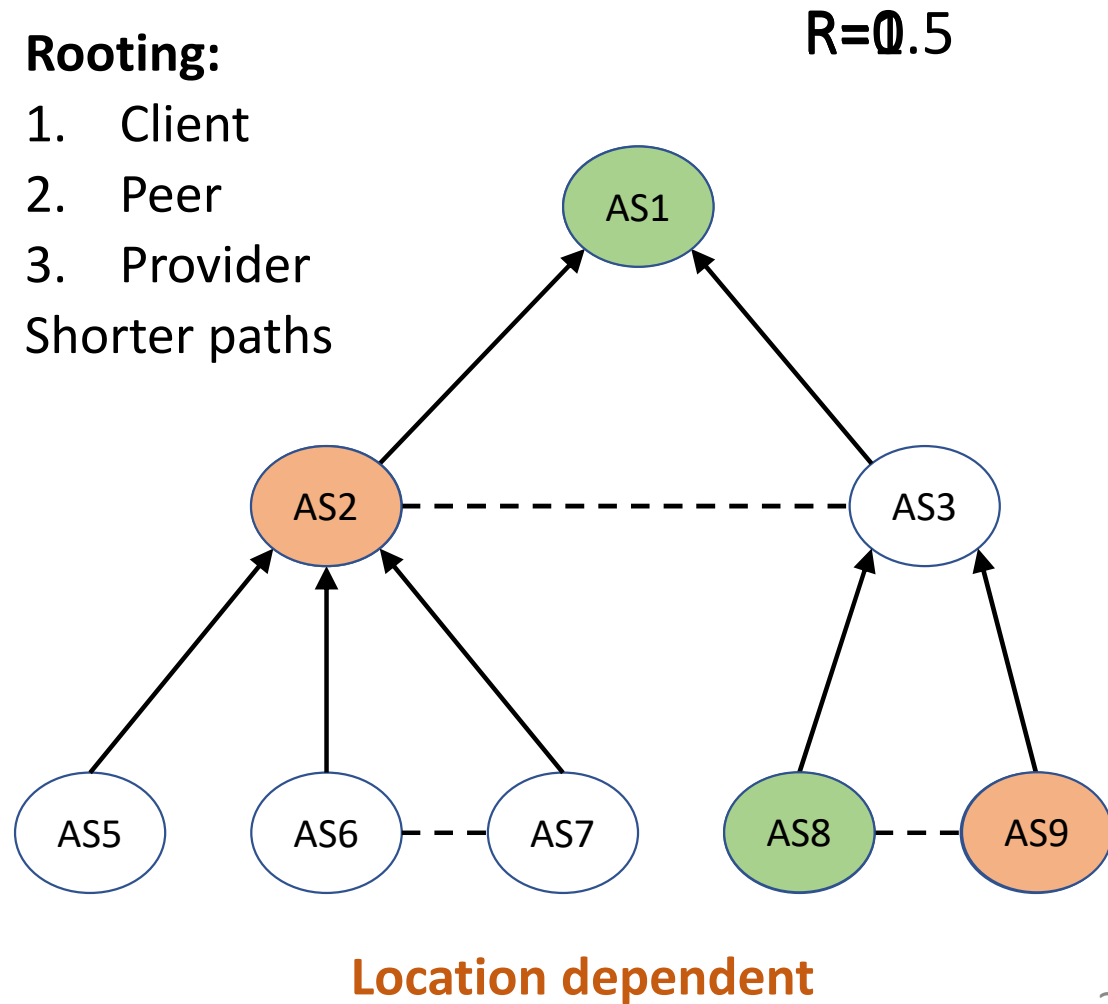
- **Privacy has become a concern**
- **Access to the Internet is censored in many countries**



- **The Tor network** - *most popular anonymization network*
  - ◆ Sender anonymity: hides the IP addresses of users
- **Problem:** Tor does not protect against global network adversary
  - ◆ Known to be vulnerable to *traffic correlation* attacks
  - ◆ Autonomous systems (ASs) apply *active routing attacks* to put themselves at both path ends
  - ◆ Alarming observations registered (WPES '04, CCS '09, CCS '13, Usenix Sec '15)

# Motivation: Counter-RAPTOR & DPSelect

- Analysis for top-93 TOR client ASs
- Performance comparable to Vanilla TOR (shadow experiments)
- **Counter-RAPTOR** (S&P '17):  $\alpha = 0.5$   
$$W_i = \alpha R_i + (1 - \alpha) \bar{B}_i$$
  - Client resilience is improved
  - No much of information leakage (mean)
- **DPSelect** (PETS'19):  
$$W_i = e^{\epsilon(\alpha(R_i)^{x_1} + (1-\alpha)(\bar{B}_i)^{x_2})}$$
  - Vulnerabilities of Counter-RAPTOR
    - Information leakage over multiple observations
    - Worst case analysis
  - Differential privacy
  - Comparable resilience



# Our Evaluation Scenario

- ***Our doubt:*** AS resilience is client-specific and easy predictable
- Potential attacker: *malicious Tor middle node*

**Do Counter-RAPTOR and DPSelect increase the vulnerability of a Tor client to a malicious middle node?**

# Our datasets

## Info about:

1. Guards ASs
2. User ASs
3. AS relationships

Description	Number	Countries	Guards	Dataset
Total number of collected ASs	57,015	230	–	–
Total number of possible user ASs	25,881	223		$D$
Total number of guard ASs	475	50	2,451	
Number of user ASs with latency	7,052	187		$D_{lat}$
Number of guard ASs with latency	333	48	2,180	

91% of IPs

89% of guard ASs

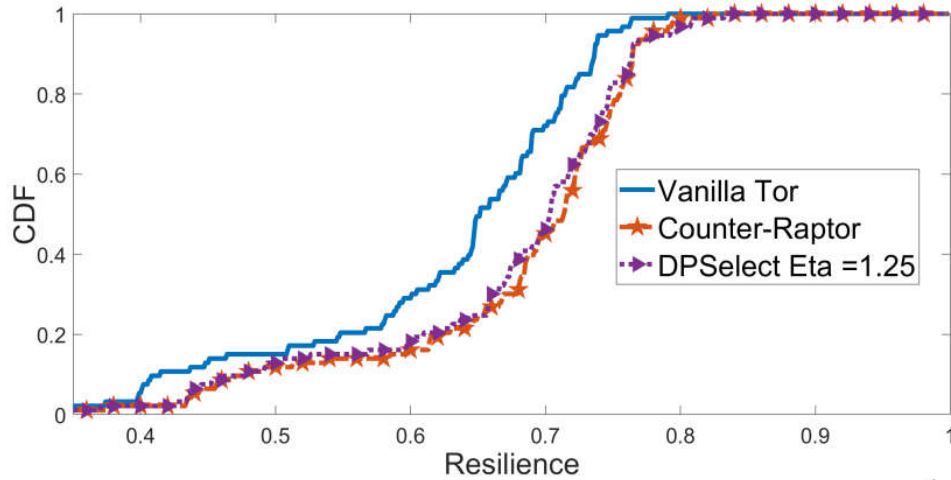
## Sources:

- CAIDA March 2017 – *ASs and relationships*
- CollecTor March 1, 2017 – *guards*
- Wacek, C., et al.: An Empirical Evaluation of Relay Selection in Tor. In: NDSS (2013) – *reduced map of the Internet including latency measurements between hosts*

**Large scale (previous works:  
top-93 TOR client ASs)**

# Our Findings (1/6)

Top-93 user ASs



$$W_i = \alpha R_i + (1 - \alpha) \bar{B}_i$$

$$R_i \in [0, 1]$$

$$\bar{B}_i = \frac{B_i}{\max_i B_i}$$

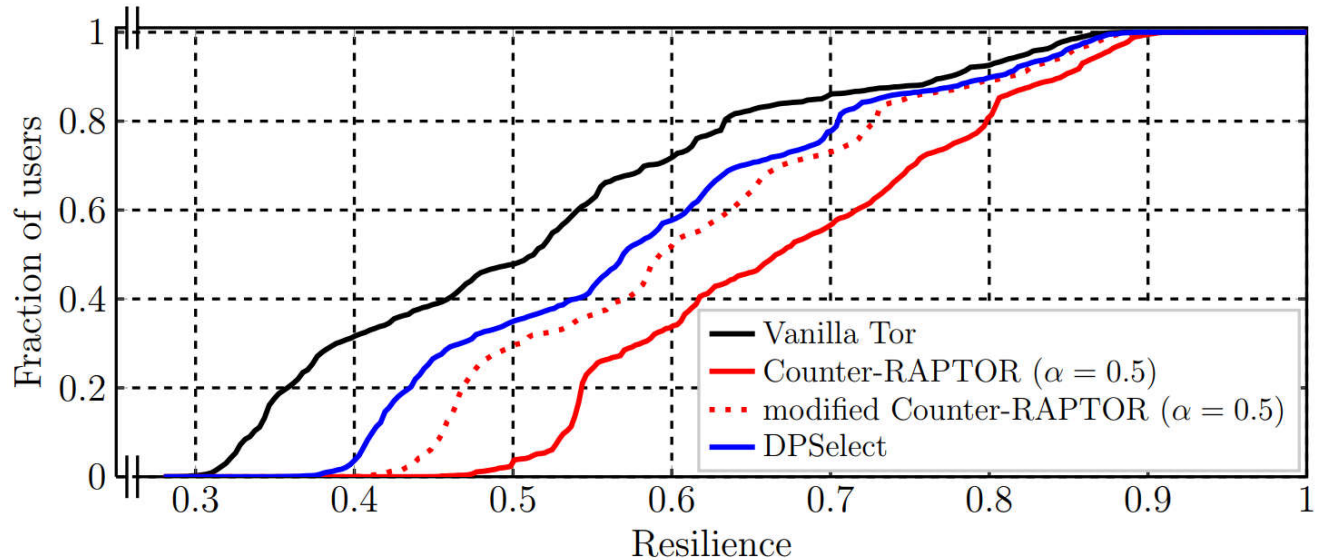
$$R_i^* = \frac{R_i}{\sum_i R_i}$$

$$B_i^* = \frac{B_i}{\sum_i B_i}$$

Hanley, Hans, et al. "DPSelect: A differential privacy based guard relay selection algorithm for tor." *Proceedings on Privacy Enhancing Technologies* 2019.2 (2019): 166-186.

DPSelect is not as good as the original Counter-RAPTOR

25881 user ASs



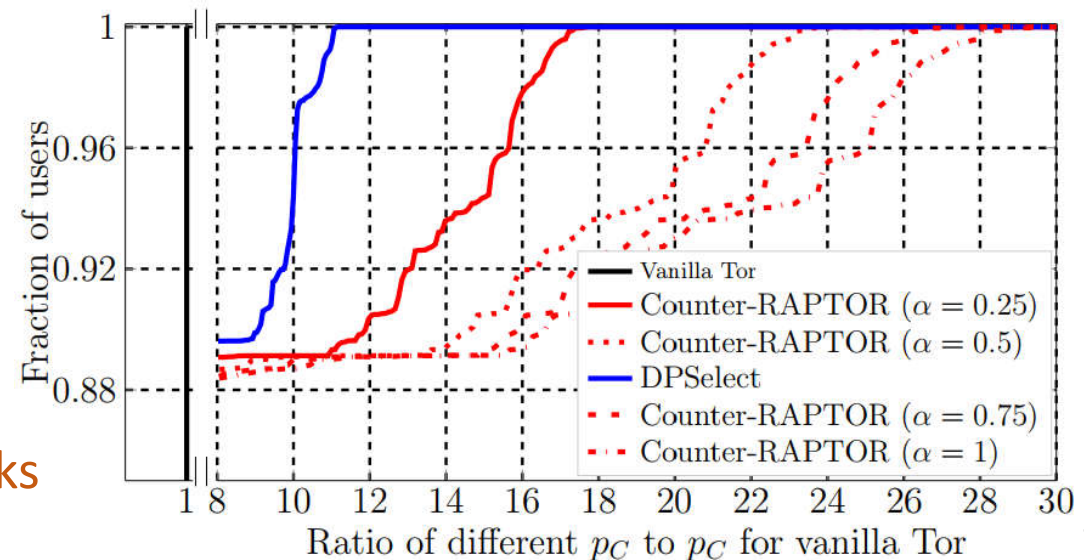
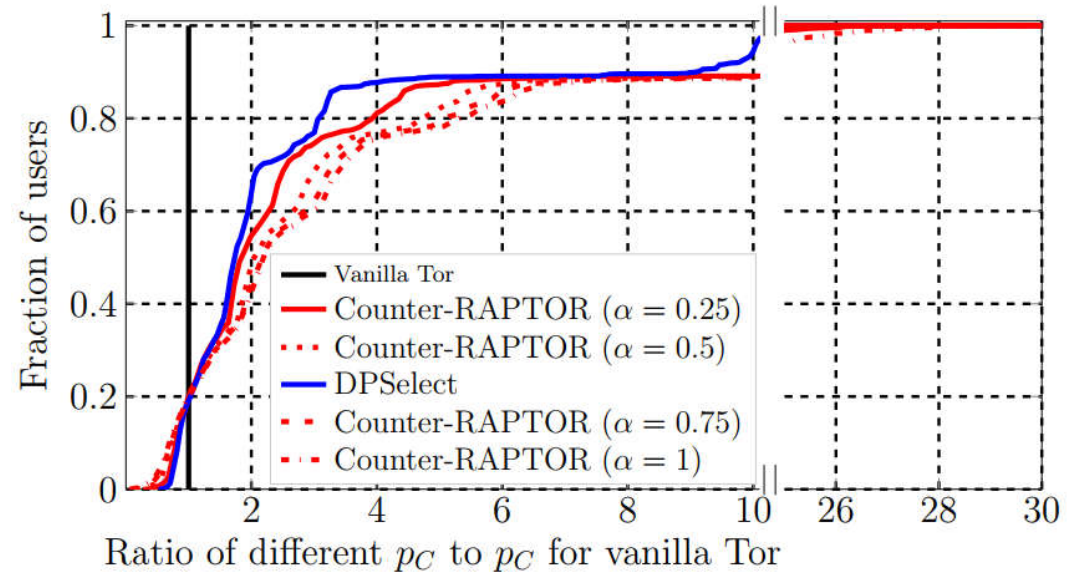
# Our Findings (2/6)

- **Geo-information leakages**

- ◆ *Hypothesis:* Counter-RAPTOR & DPSelect leak information about client location
- ◆ What about geographical position?
- ◆ Is a client more probable to choose an entry from the same country?
- ◆ *Our metric:* probability to select a guard from the same country as client

$$\frac{p_c}{p_c(\text{vanillaTor})}$$

Can be used to improve guard placement attacks



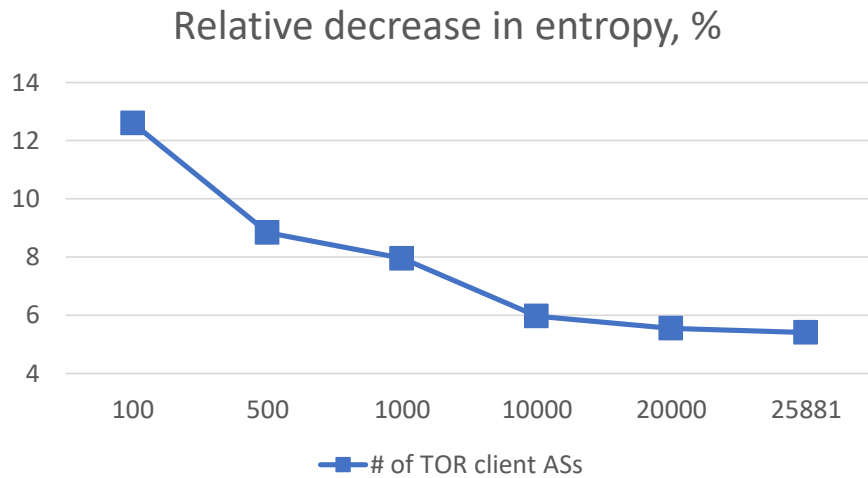
# Measuring information leakage

Counter-RAPTOR – relative decrease in entropy

*Depends on the number of client ASs*

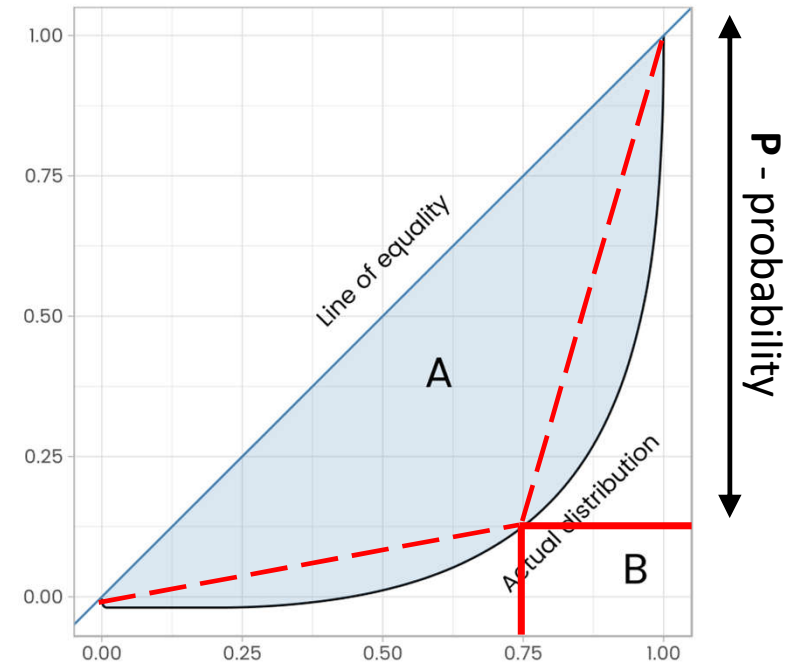


25% of probability evenly distributed between **75%** of users  
**75%** of probability evenly distributed between 25% of users



How to measure inequality

→ Gini index



$$\text{conf\_increase} = \frac{P}{F}$$

F - fraction of users  
(# of IP addresses)

We use simplified version (corresponds to 2 levels of income in economics)



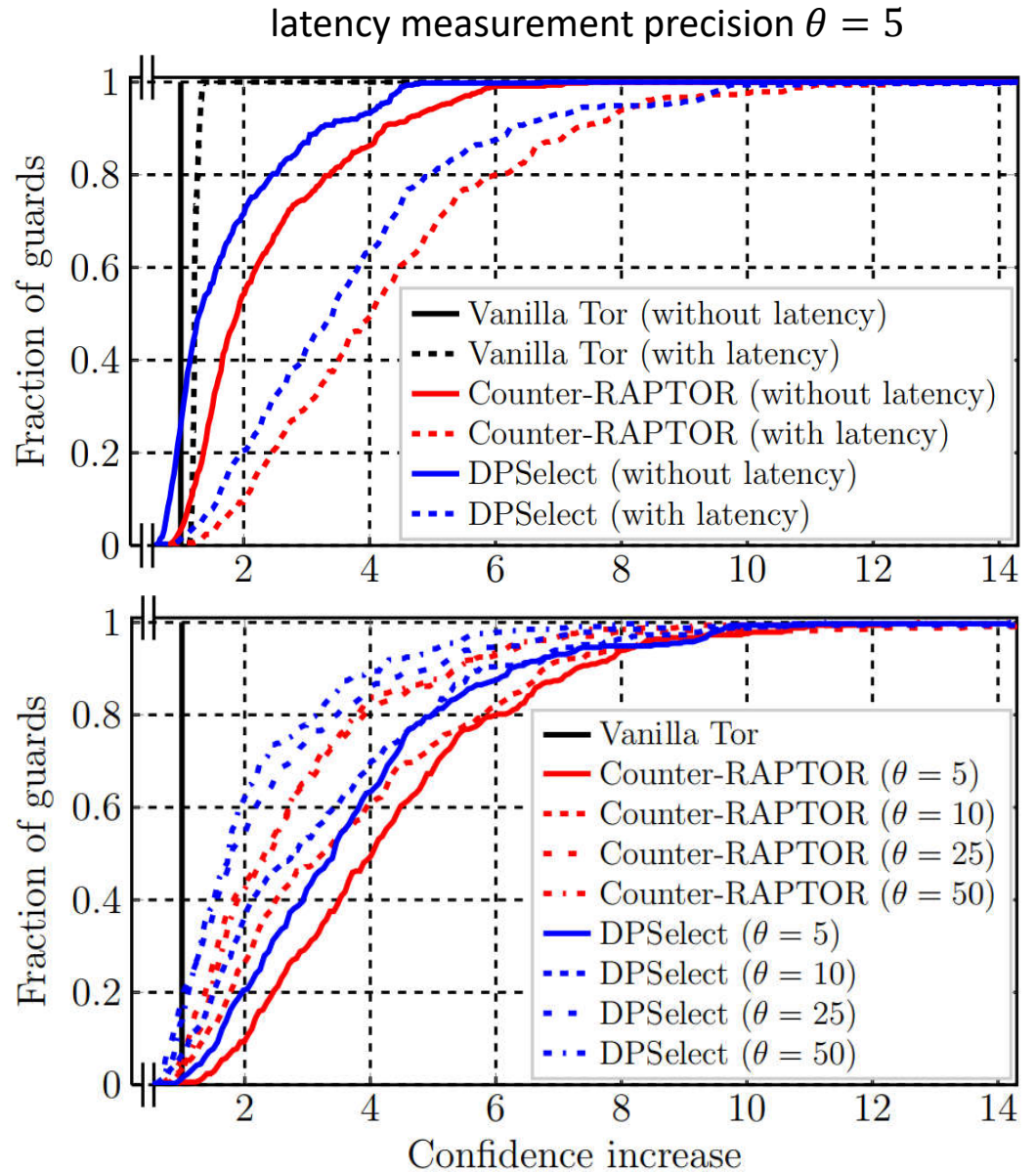
# Our Findings (3/6)

- Information gain from the position of malicious Tor middle node
- **Our metric:**

$$conf_{increase} = \frac{\text{probability}}{\text{fraction of IPs } 25\%}$$

Adding latency – simulating latency-based attacks

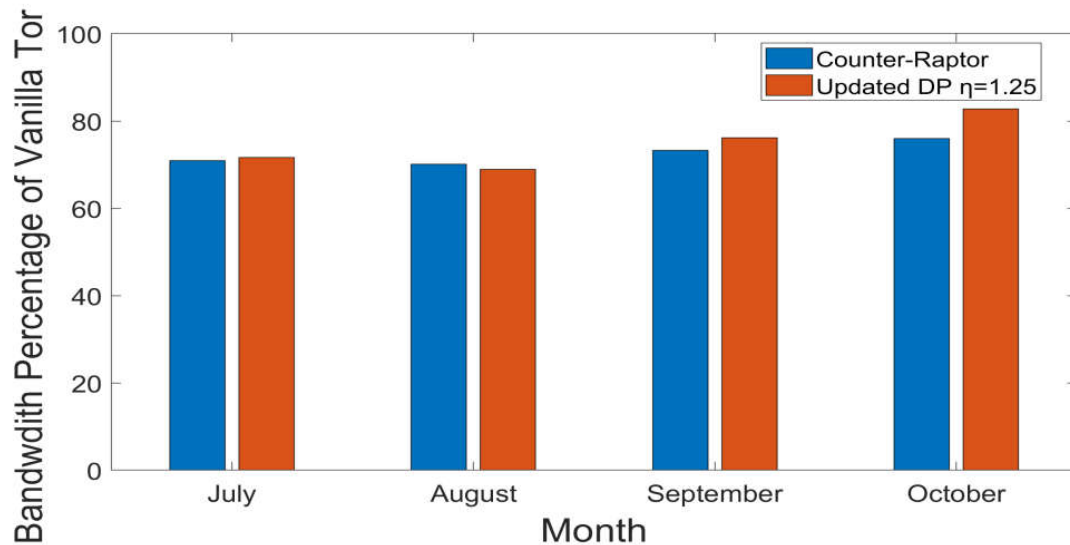
Hopper, N., et al.: How Much Anonymity Does Network Latency Leak? In: ACM CCS (2007)



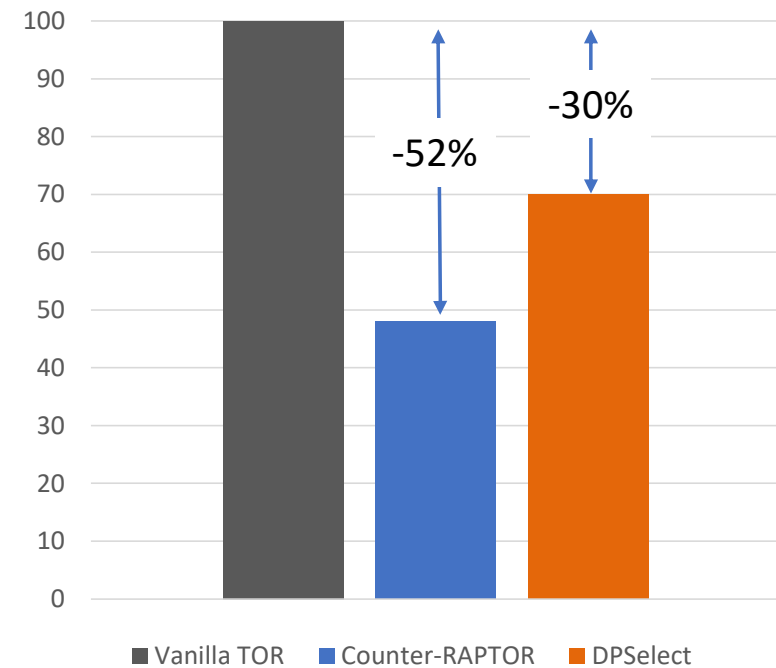
# Our Findings (4/6)

- Performance analysis
  - ◆ Average bandwidth of DPSelect in the selection of Tor entry nodes

### Top-93 user ASs



### 25881 user ASs



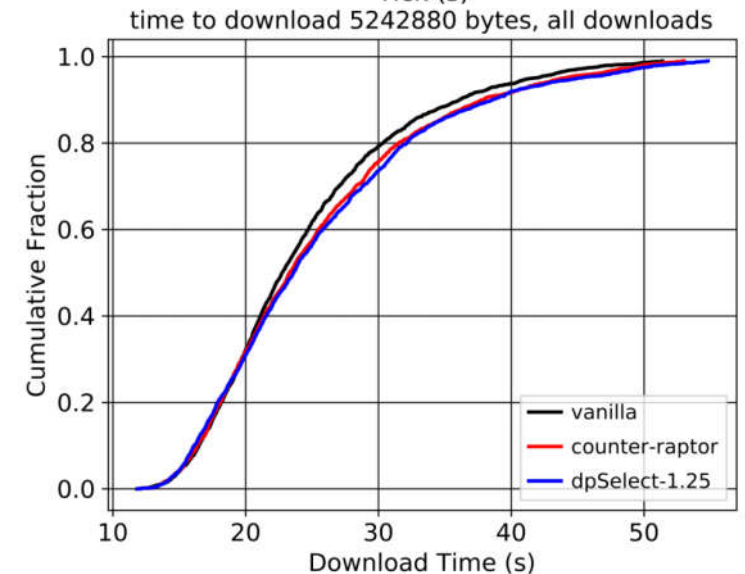
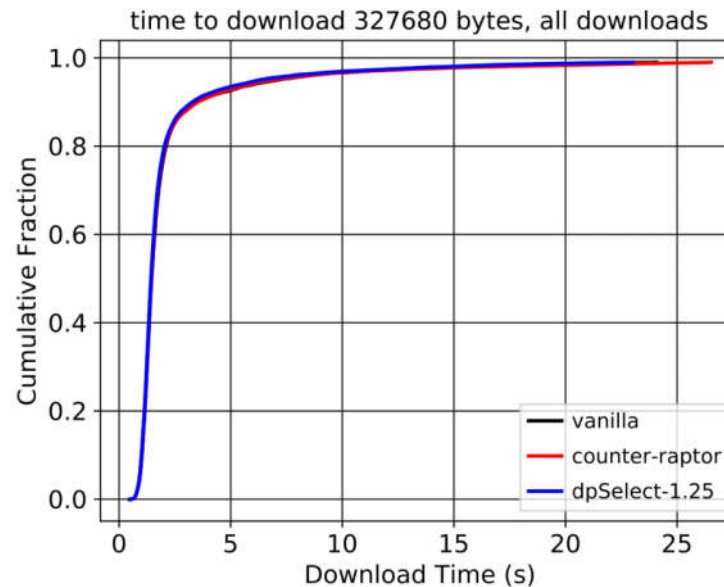
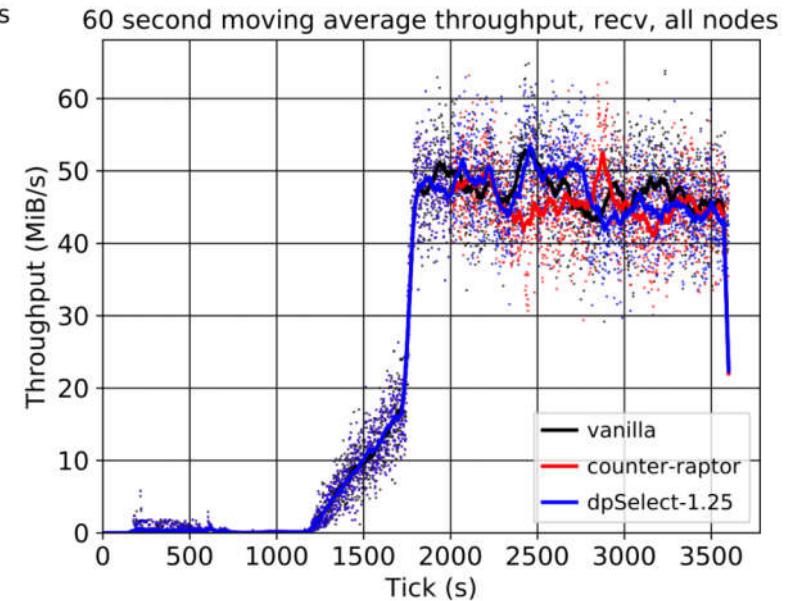
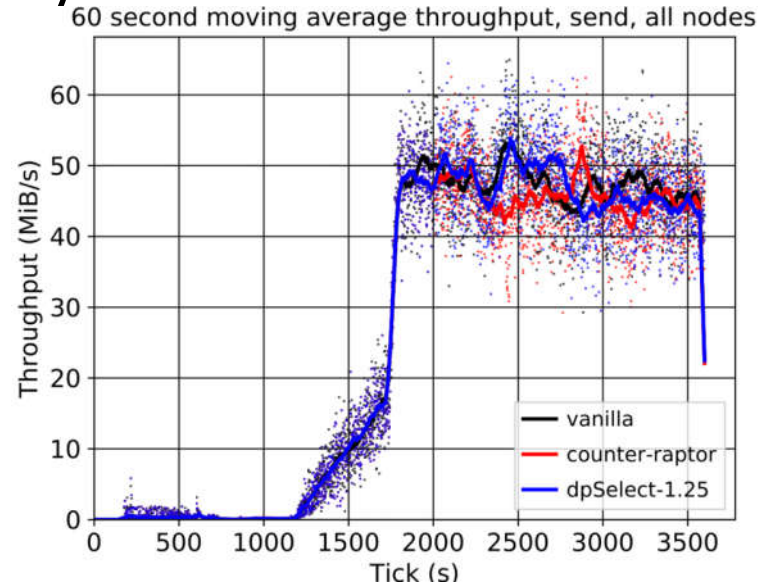
Hanley, Hans, et al. "DPSelect: A differential privacy based guard relay selection algorithm for tor." *Proceedings on Privacy Enhancing Technologies* 2019.2 (2019): 166-186.

# Performance analysis: Counter-RAPTOR & DPSelect

*Performance is similar to Vanilla TOR.*

*How can this be explained?*

Hanley, Hans, et al. "DPSelect: A differential privacy based guard relay selection algorithm for tor." *Proceedings on Privacy Enhancing Technologies* 2019.2 (2019): 166-186.

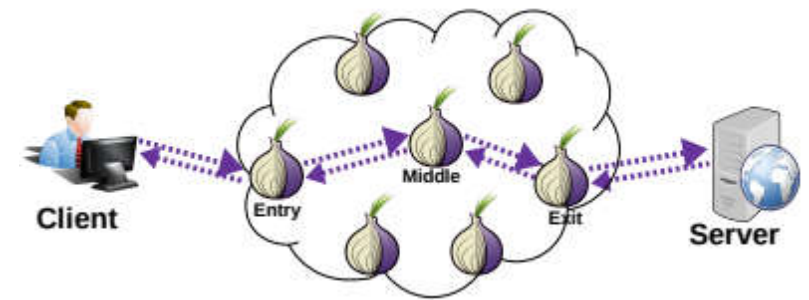


# Performance analysis: Intuition

Consensus for March 1, 2017

	$B(\text{guard})$	$B(\text{middle})$
min	1840 Kib/s	577 Kib/s
median	↑	97 Kib/s

*Only 6% of middles have greater bandwidth*

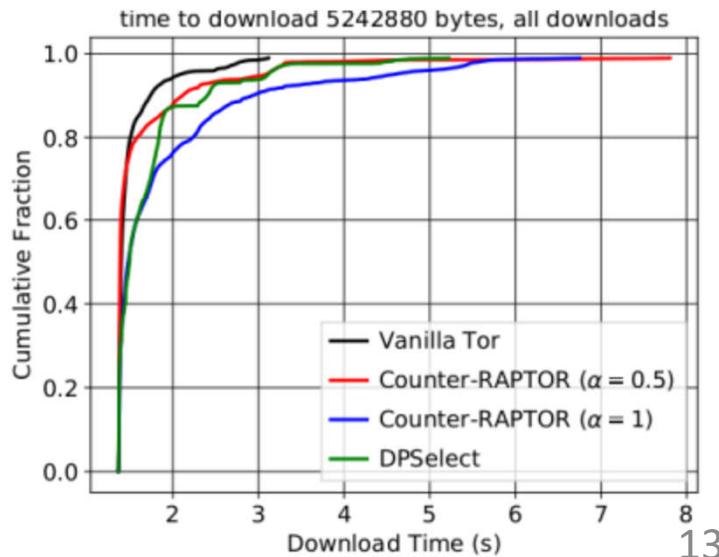
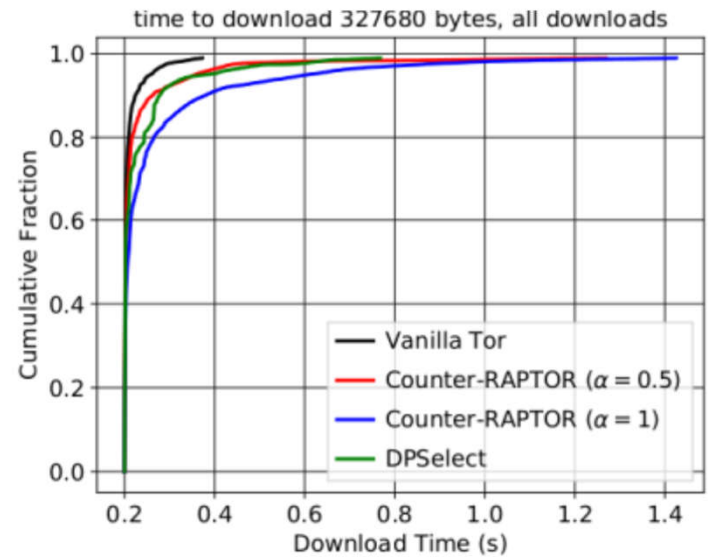
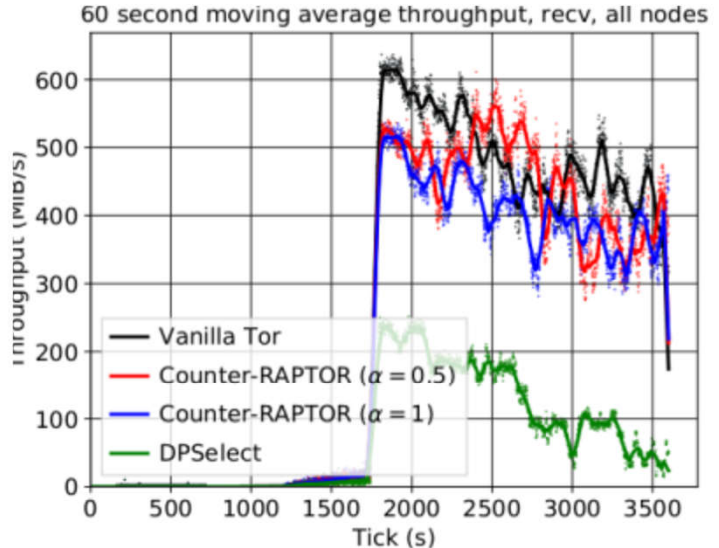
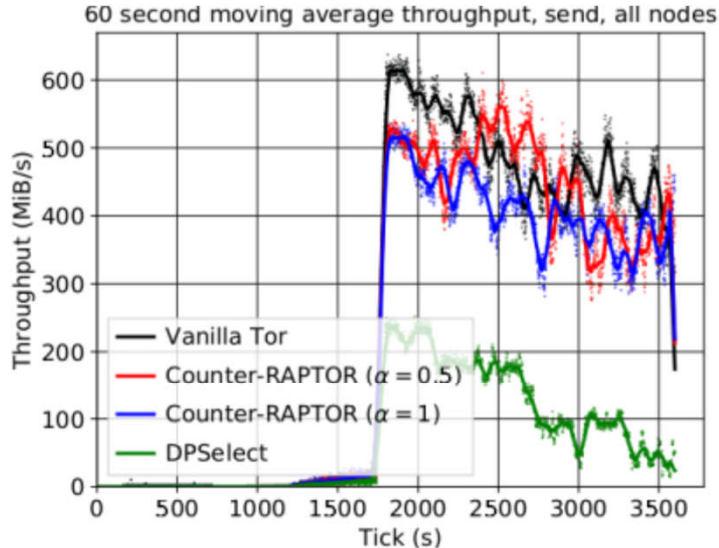


Middle relays are the bottleneck

# Our Findings (5/6)

## Experiment 1:

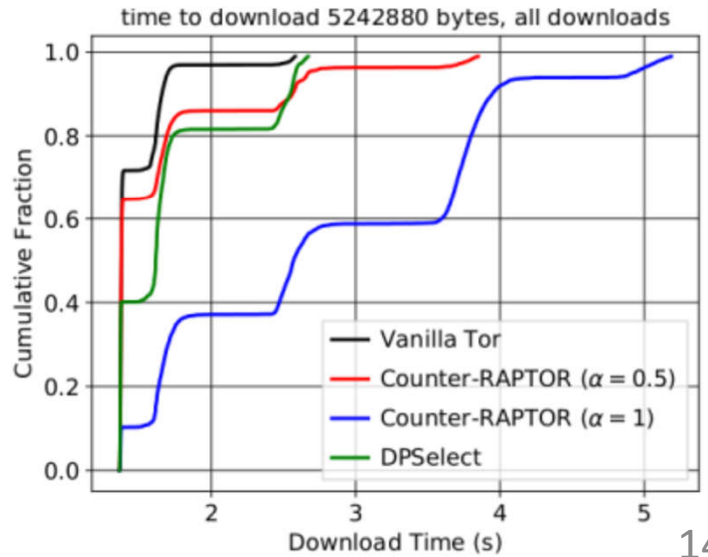
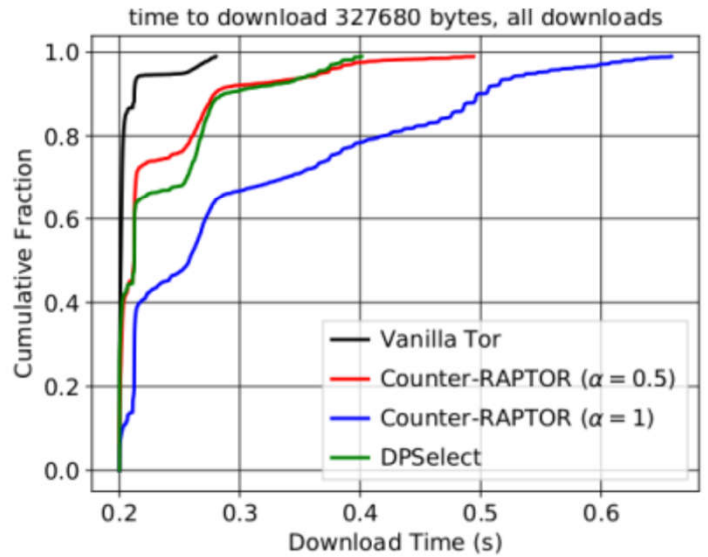
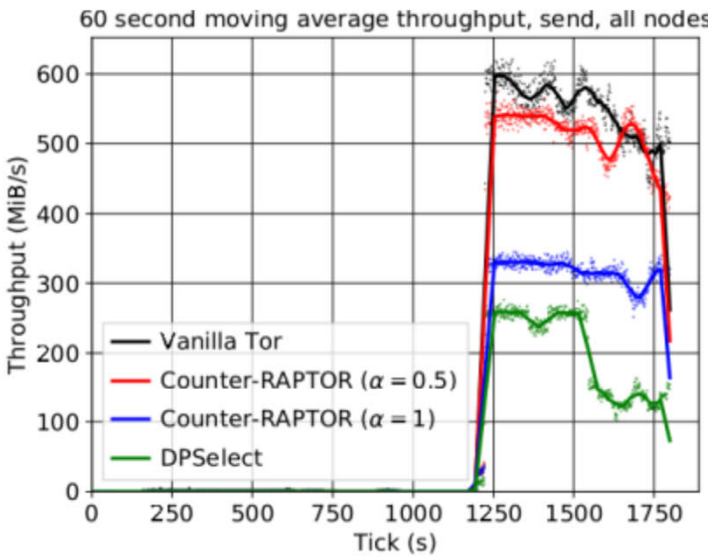
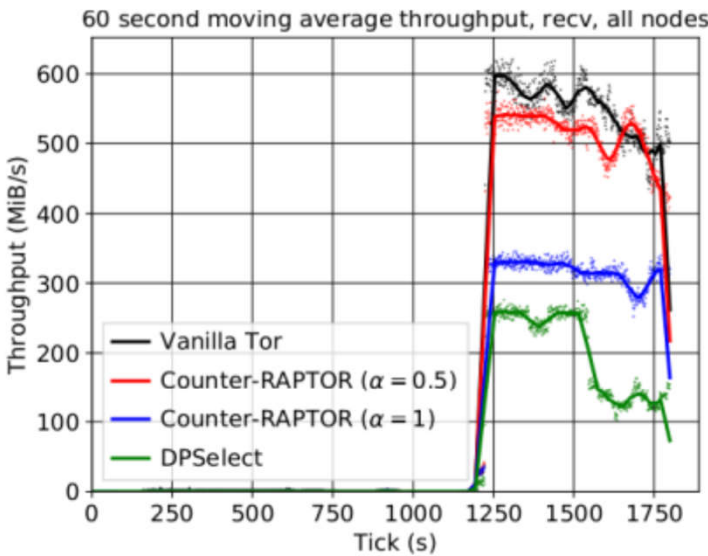
- 1.  $B(\text{middle}) \gg B(\text{guard})$
- 2.  $B(\text{exit}) \gg B(\text{guard})$
- 3. Same latency between nodes



# Our Findings (6/6)

## Experiment 2:

- 1.  $B(\text{middle}) \gg B(\text{guard})$
- 2.  $B(\text{exit}) \gg B(\text{guard})$
- 3. Same latency between nodes
- 4. All users from the same AS
- 5. 2 types of guar
  - 1. **High** performance + *low* resilience
  - 2. *Low* performance + **high** resilience



# Conclusions



## Analysis of Counter-RAPTOR & DPSelect

- DPSelect achieves only 1/3 of the claimed resilience  
→ does not protect from rooting attacks
- Both methods leak geographical information
- Analysis with regard to malicious middle OR:
  - We proposed new metric
  - Both methods empower a malicious node to fingerprint user location better
- Performance analysis
  - Degradation of average bandwidth for large scale
  - Scenarios when performance is seriously affected